

# Lecture 3

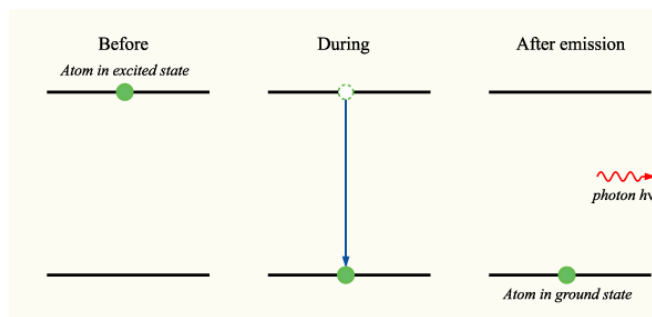
## Quantum Error Correction (QEC)

# Lecture 3

## A qubit in the real world

Consider a single qubit state  $|\psi\rangle = |1\rangle$

Due to decoherence (e.g. spontaneous emission)




$$\begin{array}{cccc}
 \underbrace{|\psi\rangle \otimes |E\rangle}_{\substack{\uparrow \\ \text{Environment}}} & \rightarrow & \sqrt{1 - p_x - p_y - p_z} |\psi\rangle \otimes |E\rangle & + \sqrt{p_x} X |\psi\rangle \otimes |E_x\rangle & + \sqrt{p_y} Y |\psi\rangle \otimes |E_y\rangle & + \sqrt{p_z} Z |\psi\rangle \otimes |E_z\rangle \\
 & & \text{No errors} & \text{Bit flip} & \text{Bit+Phase flip} & \text{Phase flip}
 \end{array}$$



How can we reset the qubit states (without destroying the qubit superposition) ?

# Quantum Error Correction

Quantum error correction: Encode a **logical** qubit in a enlarged Hilbert space, e.g., many **physical** qubits

Our first QEC code: **The three-qubit bit-flip code**

  $|\psi\rangle = \alpha |000\rangle + \beta |111\rangle$

‘Logical 0’                    ‘Logical 1’

$|0\rangle_L$                          $|1\rangle_L$

# Three-qubit bit flip code

Our first QEC code: **The three-qubit bit-flip code**

$$|\psi\rangle = \alpha |000\rangle + \beta |111\rangle$$

**The basic idea:**

One spin flips, e.g  $|\psi\rangle = \alpha |100\rangle + \beta |011\rangle$

**First step (Error syndrome): Measure  $S=Z_1Z_2, Z_2Z_3$**

→ measurement that do not destroy the superposition state, i.e. each basis state is an eigenstate of the measurement operator

→ The measurements detects the error unambiguously

# Three-qubit bit flip code

Our first QEC code: **The three-qubit bit-flip code**  $|\psi\rangle = \alpha |000\rangle + \beta |111\rangle$

**The basic idea:**

One spin flips, e.g  $|\psi\rangle = \alpha |100\rangle + \beta |011\rangle$

**First step ( Error syndrome): Measure  $S=Z_1Z_2, Z_2Z_3$**

$$\langle Z_1Z_2 \rangle = -1 \quad \langle Z_2Z_3 \rangle = 1$$

**Second step → Error recovery :  $X_1 |\psi\rangle$**

**Question:** Error syndrome/Recovery for the second and third qubit ?

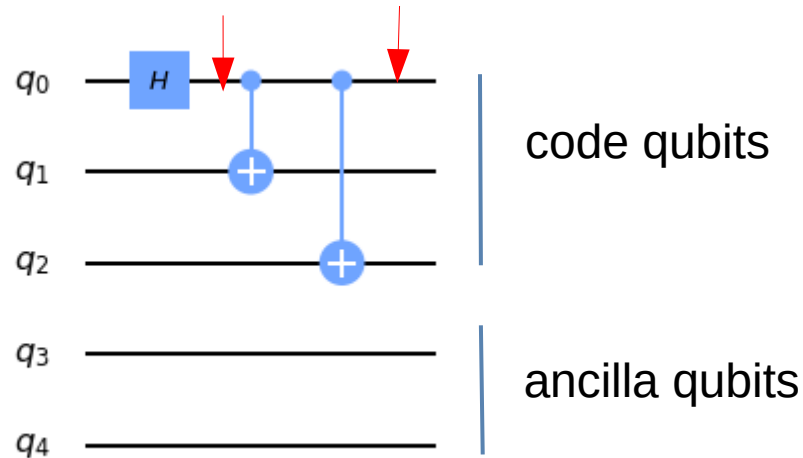
# Three-qubit bit flip code

Our first QEC code: **The three-qubit bit-flip code**

**Implementation :**

1. Create a logical qubit via two qubit entangling gates

$$|\psi\rangle = (\alpha |0\rangle + \beta |1\rangle) |00\rangle \quad |\psi\rangle = \alpha |000\rangle + \beta |111\rangle$$

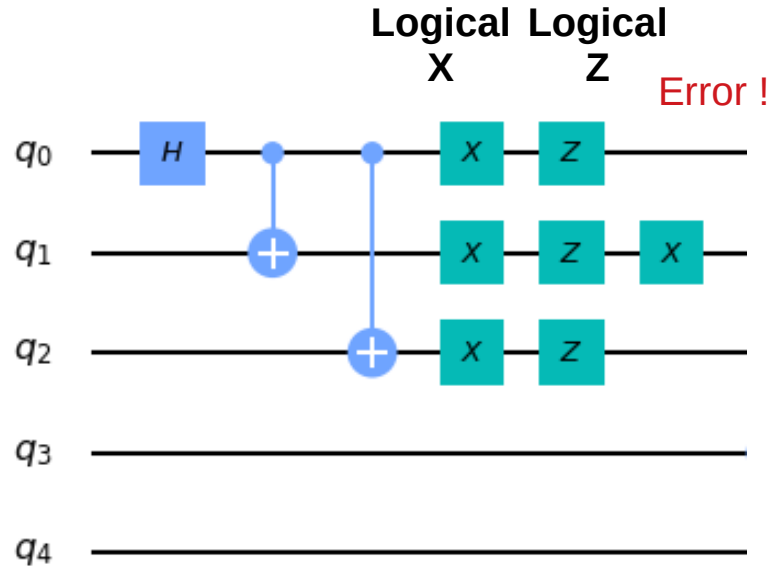


# Three-qubit bit flip code

Our first QEC code: **The three-qubit bit-flip code**

**Implementation :**

2. Evolve the logical qubit in the 'code world'

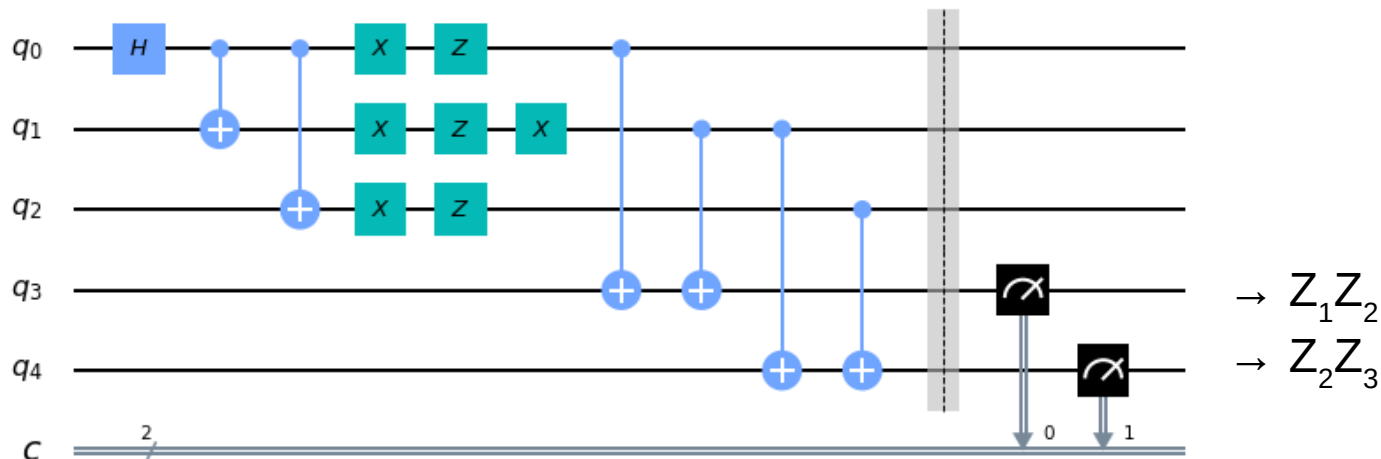


# Three-qubit bit flip code

Our first QEC code: **The three-qubit bit-flip code**

**Implementation (see TD3):**

3. Error syndrome (without destroying the qubit superposition)



4. Error recovery (easy part and optional)



# Lecture 3

Our second QEC code: **The steane code**

**Goal:** Correct arbitrary single qubit gates

$$|\psi\rangle \otimes |E\rangle \rightarrow \sqrt{1 - p_x - p_y - p_z} |\psi\rangle \otimes |E\rangle + \sqrt{p_x} X |\psi\rangle \otimes |E_x\rangle + \sqrt{p_y} Y |\psi\rangle \otimes |E_y\rangle + \sqrt{p_z} Z |\psi\rangle \otimes |E_z\rangle$$

No errorsBit flipBit+Phase flipPhase flip

For the qubit, this error decomposition **is complete** (proof in terms of Kraus representation of quantum processes)

The Steane code corrects for bit flip, phase flips, and bit+phase flips, thus for arbitrary single qubit errors.

# The Steane code

## Code world

$$|\psi\rangle = \alpha |0\rangle_L + \beta |1\rangle_L$$

$$\begin{aligned} |0\rangle_L &= |0000000\rangle + |1010101\rangle + |0110011\rangle + |1100110\rangle \\ &\quad + |0001111\rangle + |1011010\rangle + |0111100\rangle + |1101001\rangle \end{aligned}$$

$$|1\rangle_L = X_{1111111} |0\rangle_L.$$

→ 1 logical qubits = 7 physical qubits



# The Steane code

$$\begin{aligned}
 |0\rangle_L &= |0000000\rangle + |1010101\rangle + |0110011\rangle + |1100110\rangle \\
 &\quad + |0001111\rangle + |1011010\rangle + |0111100\rangle + |1101001\rangle \\
 |1\rangle_L &= X_{1111111} |0\rangle_L .
 \end{aligned}$$

Syndrome\Error	0	$X_1$	$X_2$	$X_3$	$X_4$	$X_5$	$X_6$	$X_7$	$Z_1$	$Z_2$	$Z_3$	$Z_4$	$Z_5$	$Z_6$	$Z_7$	$Y_1$	...
$Z_4 Z_5 Z_6 Z_7$	1	1	1	1	-1	-1	-1	-1	1	1	1	1	1	1	1	1	
$Z_2 Z_3 Z_6 Z_7$	1	1	-1	-1	1	1	-1	-1	1	1	1	1	1	1	1	1	
$Z_1 Z_3 Z_5 Z_7$	1	-1	1	-1	1	-1	1	-1	1	1	1	1	1	1	1	-1	
$X_4 X_5 X_6 X_7$	1	1	1	1	1	1	1	1	1	1	1	-1	-1	-1	-1	1	
$X_2 X_3 X_6 X_7$	1	1	1	1	1	1	1	1	1	-1	-1	1	1	-1	-1	1	
$X_1 X_3 X_5 X_7$	1	1	1	1	1	1	1	1	-1	1	-1	1	-1	1	-1	-1	

## Strict Conditions for QEC :

- (1) The code world is not affected by the syndrome measurement
- (2) A unique error syndrome per error

# Stabilizer formalism

## Strict Conditions for QEC :

- (1) Any logical qubit state is not affected by the syndrome measurement
- (2) A unique error syndrome per error

**These conditions can be easily checked for codes that are written in the stabilizer formalism**

## Stabilizer code $[n,k]$ (Gottesman and coworkers, 1997)

- Consider  $n$  physical qubits.
- Consider  $S$  a subgroup in the group of Pauli matrices generated by  $(n-k)$  commuting elements  $g_1, \dots, g_{n-k}$  (-Identity is excluded)
- Then  $V_S$ , the vector space stabilized by  $S$ , is of dimension  $2^k$ , i.e. we can encode  $k$  logical qubits
- The set of possible Errors  $\{E_k\}$  that can be corrected are such that for any  $j,k$ 
  - (1)  $E_j^{\text{dag}} E_k$  is in  $S$
  - Or (2)  $E_j^{\text{dag}} E_k$  anticommutes with one element of  $S$

# Stabilizer formalism

## Stabilizer code $[n,k]$

$V_S$ , the vector space stabilized by  $S$ , is of dimension  $2^k$ , i.e. can encode  $k$  logical qubits

The set of possible Errors  $E=\{E_k\}$  that can be corrected are such that for any  $j,k$

(1)  $E_j^{\text{dag}}E_k$  is in  $S$

Or (2)  $E_j^{\text{dag}}E_k$  anticommutes with one element of  $S$

**QEC recipe:** The code world is given by two orthonormal states of the vector space  $V_S$

Error Syndromes : Stabilizer measurements

**Example** **The Bit-flip code is a  $[3,1]$  stabilizer code!**

$E=\{I, X_1, X_2, X_3\}$  can be corrected by measuring the stabilizer generators  $S=\{Z_1Z_2, Z_2Z_3\}$

The logical states are  $|000\rangle$ ,  $|111\rangle$  belong to the vector space  $V_S$

**Exercise** Prove that the Steane code is a  $[7,1]$  stabilizer code.

# The challenge of fault-tolerant quantum computing

A single QEC code does not protect against **any** error. (ex error  $X_1X_2$  in the bit flip code..)

However, we may concatenate/combine several QEC codes to fight for error propagation provided the error probability per gate is below a certain threshold [Michael Ben-Or and Dorit Aharonov]

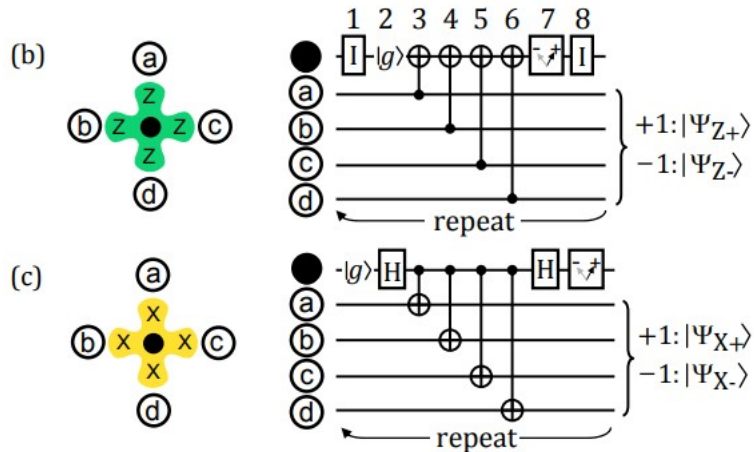
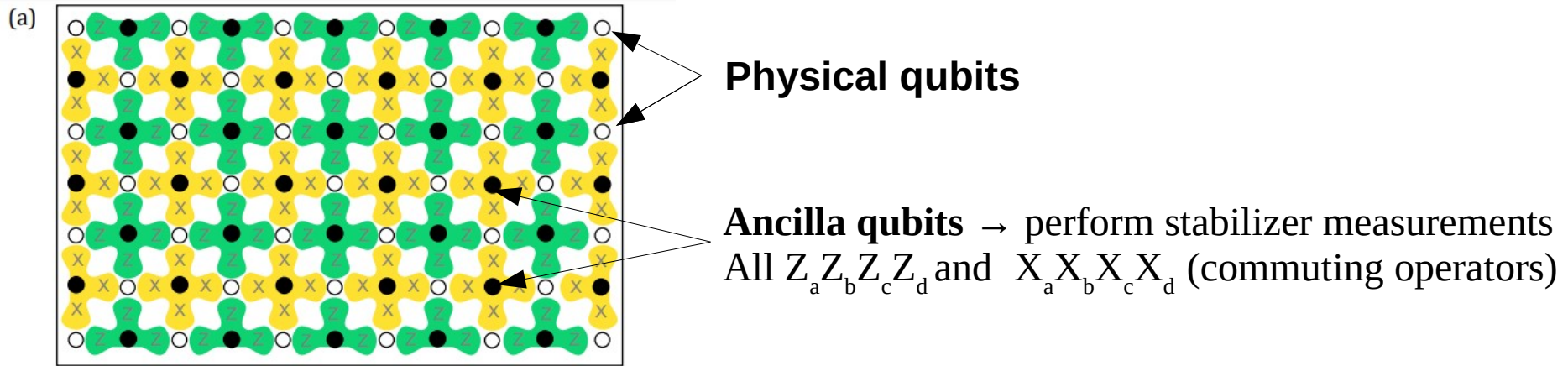
**Fault-tolerant quantum computing = A significant technological challenge.**

Estimation : ~10 000 physical qubits per logical qubit for the surface code



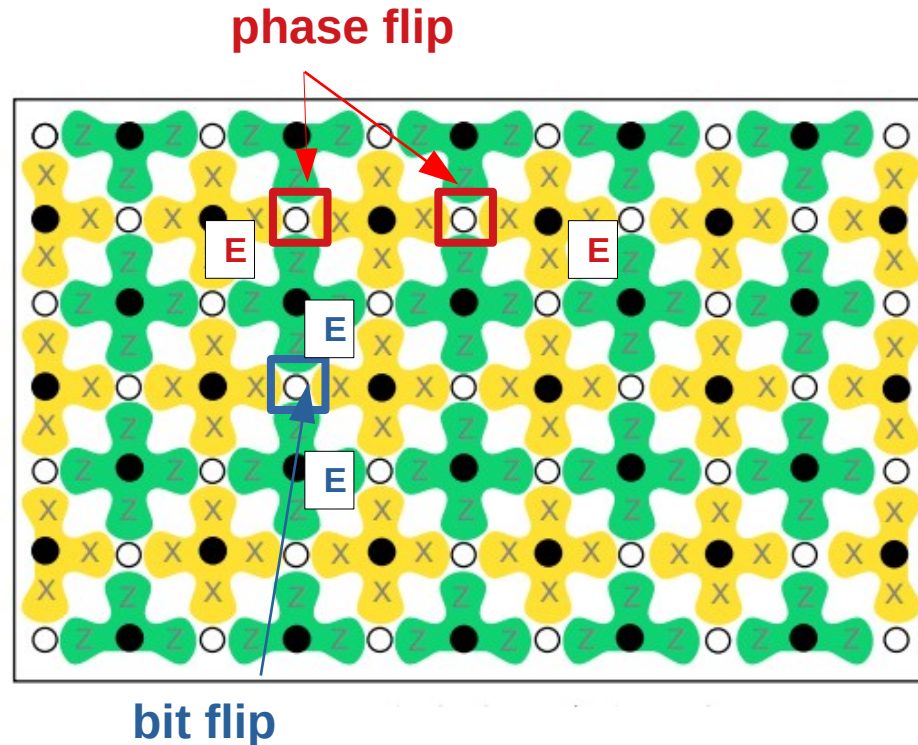
*J. Preskill « I've already emphasized repeatedly that it will probably be a long time before we have fault-tolerant quantum computers solving hard problems. »*

# Towards fault-tolerance with the surface code



# Towards fault-tolerance with the surface code

Error syndrome = pattern of stabilization results

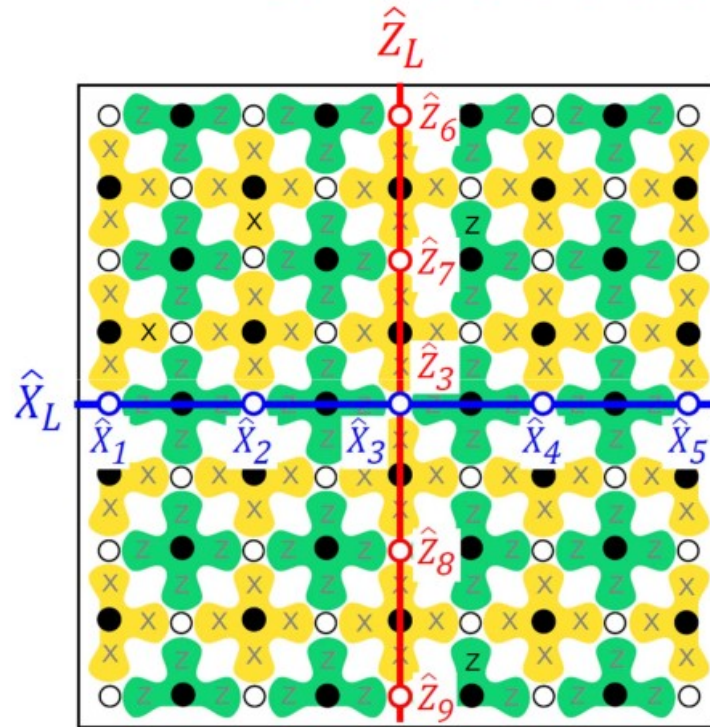




# Towards fault-tolerance with the surface code

## Initialization and operation

- 1) Measure all stabilizers  
→ Creates a logical state  $|\psi\rangle$
- 2) Apply logical operations  
 $X_L$   $Z_L$  (Preserve the code world)
- 3) Check for errors by stabilizer measurements

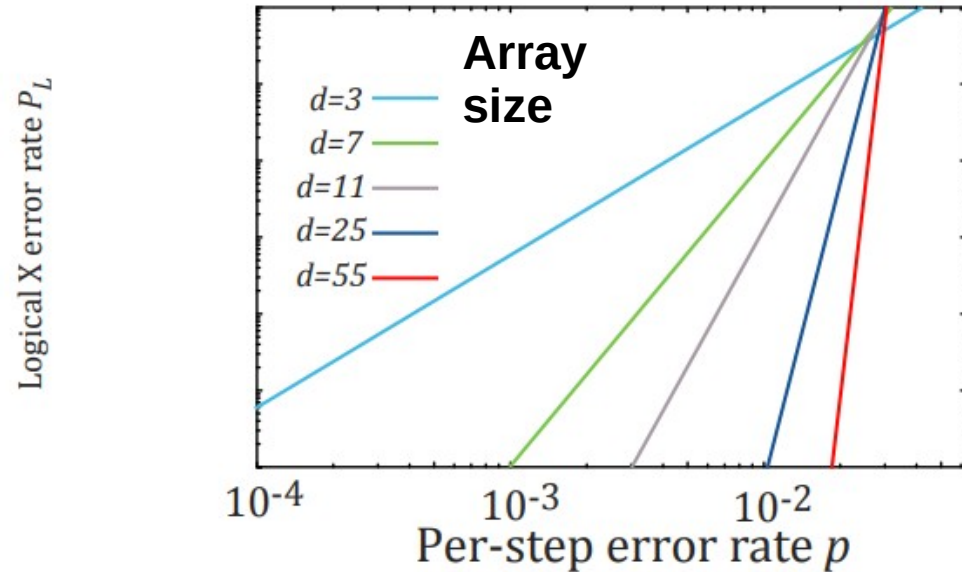
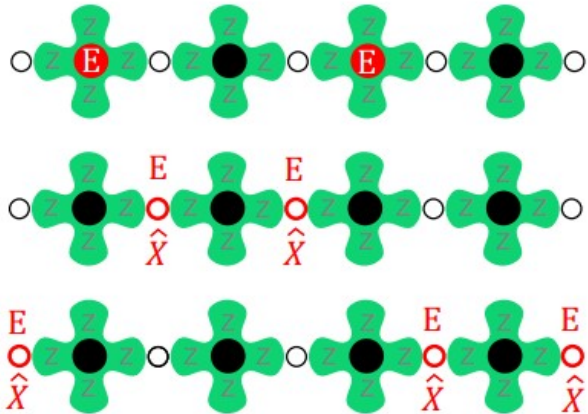


# Towards fault-tolerance with the surface code

## Robustness of the logical qubit (TD3)

QEC fails when the numbers of errors  $\sim$  system size  $d/2$

**Example** : One error syndrome for two possible errors

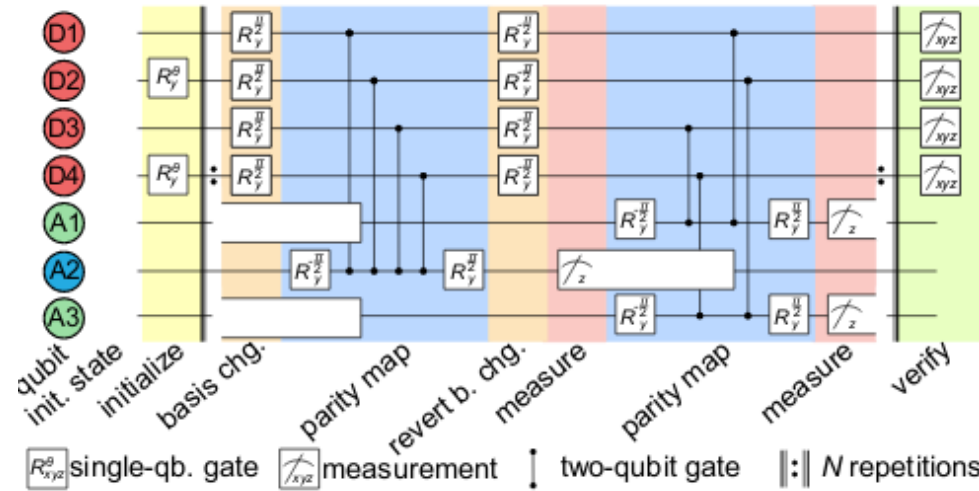
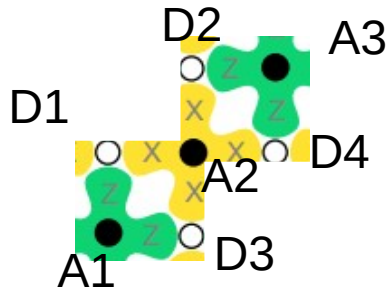


Fault-tolerance can be achieved by increasing the array size

# Towards fault-tolerance with the surface code

## Experimental implementation of a minimal surface code

<https://arxiv.org/pdf/1912.09410.pdf>



Stabilizer :  $Z_{D1}Z_{D3}$ ,  $Z_{D2}Z_{D4}$ , and  $X_{D1}X_{D2}X_{D3}X_{D4}$

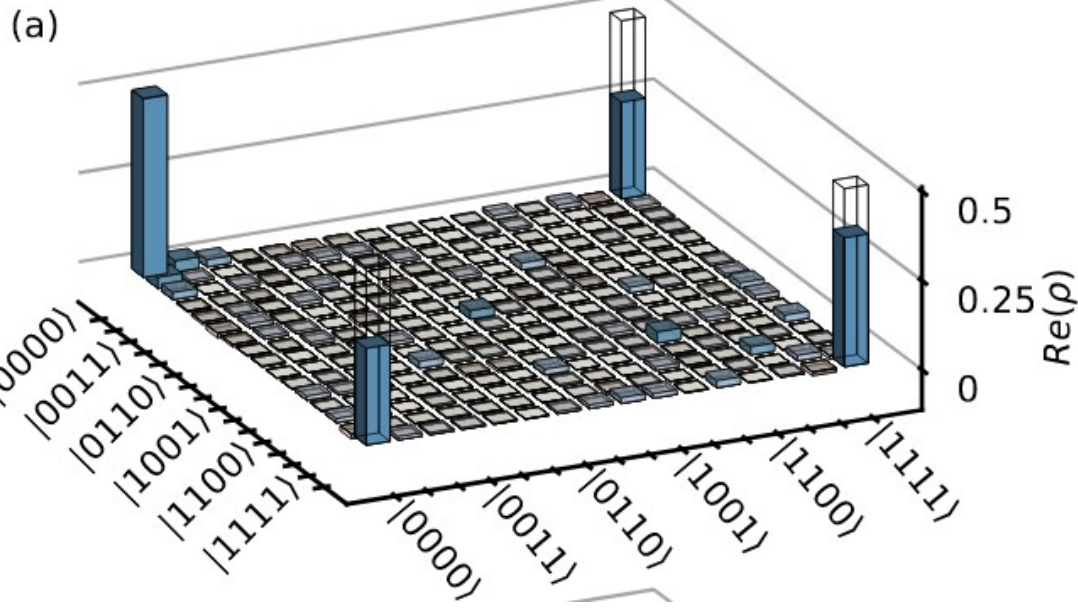
$$|0\rangle_L = \frac{1}{\sqrt{2}} (|0000\rangle + |1111\rangle)$$

$$|1\rangle_L = \frac{1}{\sqrt{2}} (|0101\rangle + |1010\rangle)$$

# Towards fault-tolerance with the surface code

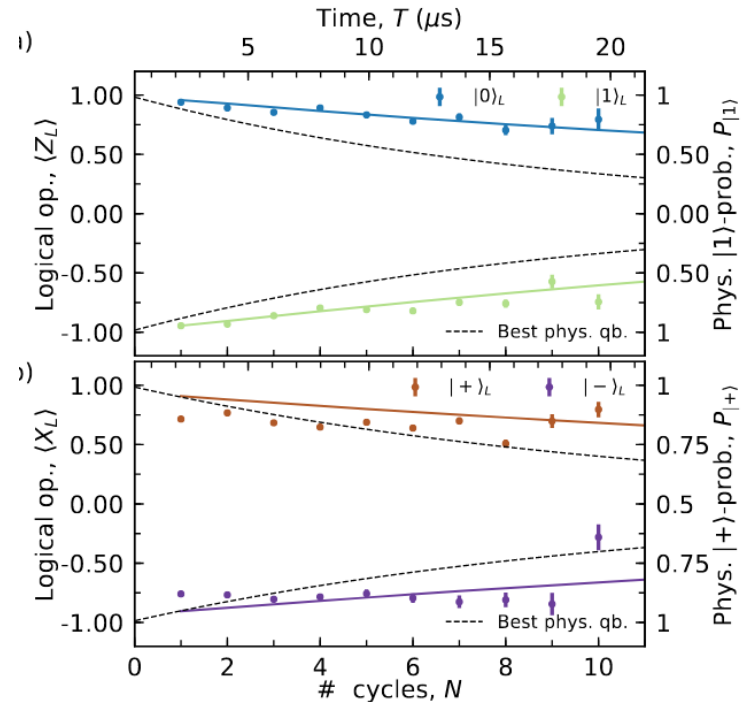
## Preparation of the $0_L$ logical state

- start from  $0^4$
- measure all ancillas in 0



## Repeated error correction

(postprocessing on the no-error signal)



# Consequence : we need more physical qubits

## Scaling IBM Quantum technology



IBM Q System One (Released)

(In development)

Next family of IBM Quantum systems

2019

2020

2021

2022

2023

and beyond

27 qubits

65 qubits

127 qubits

433 qubits

1,121 qubits

Path to 1 million qubits

*Falcon*

*Hummingbird*

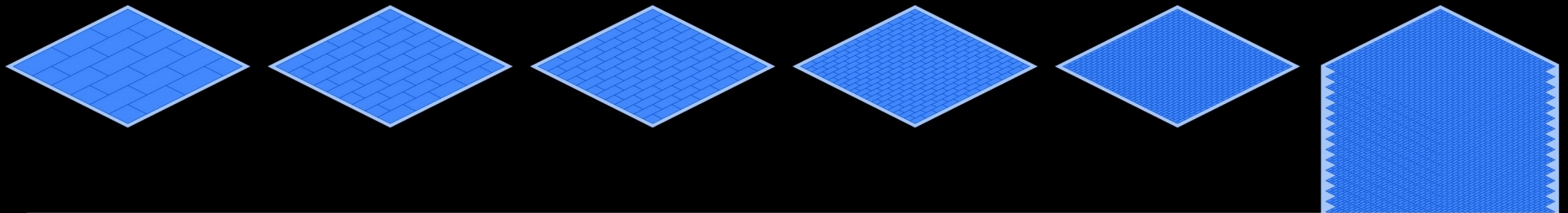
*Eagle*

*Osprey*

*Condor*

and beyond

*Large scale systems*



Key advancement

Key advancement

Key advancement

Key advancement

Key advancement

Key advancement

Optimized lattice

Scalable readout

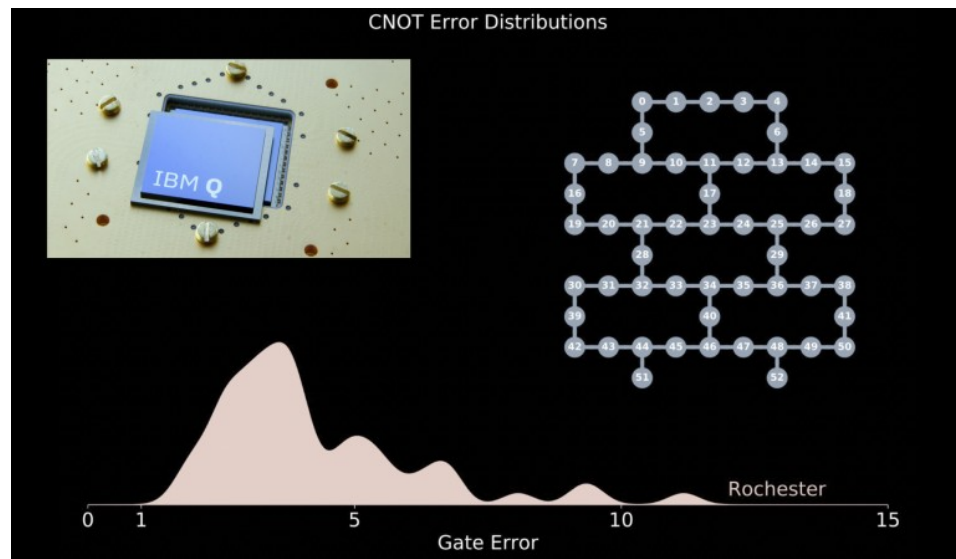
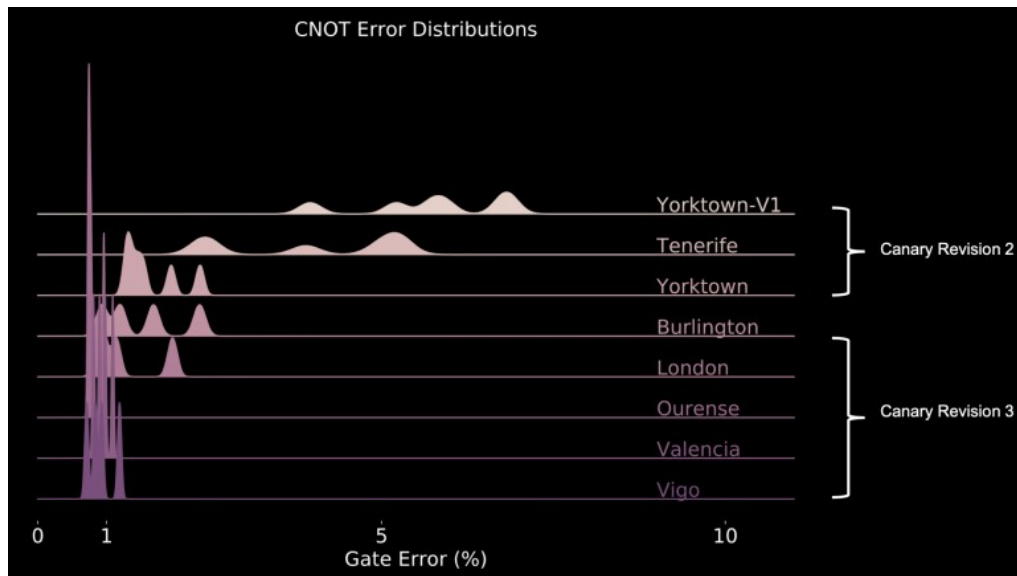
Novel packaging and controls

Miniaturization of components

Integration

Build new infrastructure,  
quantum error correction

# Consequence : we need more physical qubits





# Lecture 3 : Summary

QEC is a hot topic in modern quantum computing and **a key challenge** for quantum technologies

**Fault-tolerant quantum computing** requires many physical qubit per logical qubit

**Surface codes** implement QEC as 2D spin lattice models, which we also encounter in quantum simulation : Lecture 4

